

1. For each of the following, translate the given argument into symbolic form. Then determine whether or not the given argument is valid by either writing a 2-column proof for the argument or by finding a counterexample.

- (a) If I work hard every day then I will get a promotion at work.
 If I do not get a promotion at work then I will not be able to afford my house payment.
 I can afford my house payment.

 Therefore I work hard every day.

We assign variables: w : I work hard every day. p : I get a promotion at work. a : I can afford my house payment.

The the symbolic form of this argument is:

$$\begin{array}{l} w \rightarrow p \\ \neg p \rightarrow \neg a \\ \hline a \\ \hline \therefore w \end{array}$$

This argument is not valid. Notice that we can use the final two premises along with the Law of Contraposition to conclude p : this person got a promotion. However, concluding that this person worked hard in order to get a promotion, based on the first premise would be the Fallacy of the Converse. We know that working hard every day leads to getting a promotion, but we cannot assume that since they got a promotion, they must have worked hard every day. The problem truth values are when w is False, p is True, and a is True.

- (b) If gas is expensive and parking is inconvenient then I will take the bus to school.
 If I take the bus to school then I will not be able to take a night class.
 I am taking a night class.

 Therefore gas is not too expensive or parking is convenient.

We begin by assigning variables as follows: g : gas is expensive. p : parking is inconvenient. b : I take the bus to school. n : I take a night class.

Then the symbolic form of this argument is:

$$\begin{array}{l} (g \wedge p) \rightarrow b \\ b \rightarrow \neg n \\ \hline n \\ \hline \therefore \neg g \vee \neg p \end{array}$$

We now construct a 2-column proof as follows:

Statement	Reason
1. $b \rightarrow \neg n$	Premise
2. n	Premise
3. $\neg b$	Law of Contraposition (Modus Tollens) (1, 2)
4. $(g \wedge p) \rightarrow b$	Premise
5. $\neg(g \wedge p)$	Law of Contraposition (Modus Tollens) (3,4)
6. $\neg g \vee \neg p$	De Morgan's Law (5)

Hence this argument is valid.

2. Determine whether or not the given argument is valid by translating it into symbolic form and identifying its form as a known valid argument form or a known fallacy.

- (a) If I want to go out on Saturday night then I need to study for my exam during the afternoon.
 I do not want to go out on Saturday night.

 Therefore I do not need to study for my exam during the afternoon.

We assign variables: O : I want to go out on Saturday night. S : I need to study for my exam during the afternoon.

The the symbolic form of this argument is:

$$\begin{array}{l} O \rightarrow S \\ \neg O \\ \hline \therefore \neg S \end{array}$$

The form of this argument is the Fallacy of Denying the Hypothesis, so it is not valid.

- If I study for my exam Saturday during the afternoon then I will go out on Saturday night.
- (b) $\frac{\text{I do not go out on Saturday night.}}{\text{Therefore I did not study for my exam Saturday afternoon.}}$
- We assign variables: O : I will go out on Saturday night. S : I study for my exam Saturday during the afternoon.
- The the symbolic form of this argument is:

$$\frac{\begin{array}{l} S \rightarrow O \\ \neg O \end{array}}{\therefore \neg S}$$

The form of this argument is Modus Tollens, so it is valid.

- If I do not study for my exam then I will not get a good grade on it.
- (c) $\frac{\text{I got a good grade on my exam.}}{\text{Therefore I studied for my exam.}}$

We assign variables: S : I study for my exam. G : I get a good grade on my exam.

The the symbolic form of this argument is:

$$\frac{\begin{array}{l} \neg S \rightarrow \neg G \\ G \end{array}}{\therefore S}$$

The form of this argument is Modus Tollens, so it is valid.

3. Prove that each of the following arguments are valid by constructing a 2-column proof.

- If you are unhappy with the results of the election then the candidate you voted for did not win.
- If the candidate you voted for did not win, then she will run again next election.
- (a) $\frac{\text{The candidate you voted for is not running next election}}{\text{Therefore you are happy with the election results}}$

We begin by translating this argument into symbolic form. Defining the variables H : You are happy with the results of the election. W : The candidate you voted for won. R : The candidate you voted for is running in the next election, we have the following argument:

- (b) $\frac{\begin{array}{l} \neg H \rightarrow \neg W \\ \neg W \rightarrow R \\ \neg R \end{array}}{\therefore H}$

Then we can construct the following proof:

Statement	Reason
$\neg H \rightarrow \neg W$	Premise
$\neg W \rightarrow R$	Premise
$\neg H \rightarrow R$	Hypothetical Syllogism
$\neg R$	Premise
H	Modus Tollens

Hence this argument is valid.

- Everyone who is brave and intelligent majors in mathematics.
- Tony is not a math major.
- (c) $\frac{\text{Tony is brave.}}{\text{Therefore Tony is not intelligent.}}$

We begin by translating this argument into symbolic form. Let x represent people. We then defining the predicates $B(x)$: Person x is brave. $I(x)$: Person x is intelligent. $M(x)$: Person x is a math major. Then we have the following argument:

- (d) $\frac{\begin{array}{l} \forall x [(B(x) \wedge I(x)) \rightarrow M(x)] \\ \neg M(\text{Tony}) \\ B(\text{Tony}) \end{array}}{\therefore \neg I(\text{Tony})}$

Then we can construct the following proof:

Statement	Reason
$\forall x [(B(x) \wedge I(x)) \rightarrow M(x)]$	Premise
$(B(\text{Tony}) \wedge I(\text{Tony})) \rightarrow M(\text{Tony})$	Universal Instantiation
$\neg M(\text{Tony})$	Premise
$\neg (B(\text{Tony}) \wedge I(\text{Tony}))$	Modus Tollens
$\neg B(\text{Tony}) \vee \neg I(\text{Tony})$	DeMorgan's Law
$B(\text{Tony})$	Premise
$\neg I(\text{Tony})$	Disjunctive Syllogism

Hence this argument is valid.

4. Prove that the product of three odd numbers is odd.

Proof:

We will begin by proving the following Lemma:

Lemma: The product of two odd numbers is odd.

Proof: Let m and n be odd numbers. Then, by definition, there is an integer k such that $m = 2k + 1$ and there is an integer l such that $n = 2l + 1$.

Then $mn = (2k + 1)(2l + 1) = 4kl + 2k + 2l + 1 = 2(2kl + k + l) + 1$, which, since $2kl + k + l$ is an integer, shows that mn is odd. \square

From this, we can prove the original theorem as follows:

Let m, n and p be odd integers. Using the lemma, mn is odd. But then $mnp = (mn)p$ is the product of two odd integers, so again using the lemma, mnp is odd.

5. Prove that if n is an integer and that $n^2 + 11$ is even, then n is odd.

Proof:

We will proceed using Proof by Contraposition.

Suppose that n is even. Then, by definition, $n = 2k$ for some integer k . Therefore, $n^2 + 11 = (2k)^2 + 11 = 4k^2 + 10 + 1 = 2(2k^2 + 5) + 1$.

Thus $n^2 + 11$ is odd.

We have shown that if n is odd, then $n^2 + 11$ is odd, hence the contrapositive of this statement is true.

Therefore, if n is an integer and $n^2 + 11$ is even, then n is odd. \square

6. Prove that $m^2 = n^2$ if and only if $m = n$ or $m = -n$.

“ \Rightarrow ”

We begin by proving the backward direction of this biconditional statement. Suppose that either $m = n$ or $m = -n$.

If $m = n$, then $m^2 = n^2$ follows using basic principles of Algebra.

Similarly, if $m = -n$, then $m^2 = (-n)^2 = n^2$ follows using basic principles of Algebra.

“ \Leftarrow ”

To prove the backward direction of this biconditional statement, suppose that $m^2 = n^2$.

Notice that the square of any real number is non-negative, so we can take the square root of each side of this equation in order to obtain:

$$\sqrt{m^2} = \sqrt{n^2}, \text{ or } |m| = |n|.$$

Thus, using basic principles of algebra, either $m = n$ or $m = -n$. \square

7. Show that the following three statements are all equivalent:

(i) x is rational

(ii) $\frac{x}{2}$ is rational

(iii) $x + 1$ is rational.

Proof: To show that all of these statements are equivalent, we will show $(i) \rightarrow (ii) \rightarrow (iii) \rightarrow (i)$.

“(i) \rightarrow (ii)”

Assume that x is rational. Then, by definition, $x = \frac{a}{b}$ for integers a, b with $b \neq 0$.

But then $\frac{x}{2} = \frac{a}{2b}$, so $\frac{x}{2}$ is also rational.

“(ii) \rightarrow (iii)”

Assume that $\frac{x}{2}$ is rational. Then, by definition, $\frac{x}{2} = \frac{a}{b}$ for integers a, b with $b \neq 0$.

But then $2 \cdot \frac{x}{2} + 1 = x + 1 = \frac{2a}{b} + 1 = \frac{2a+b}{b}$, so $x + 1$ is also rational.

“(iii) \rightarrow (i)”

Assume that $x + 1$ is rational. Then, by definition, $x + 1 = \frac{a}{b}$ for integers a, b with $b \neq 0$.

But then $x = (x + 1) - 1 = \frac{a}{b} - 1 = \frac{a-b}{b}$, so x is also rational. \square

8. Prove or Disprove: The sum of two irrational numbers is irrational.

This theorem is false. Here is a counterexample: as we proved in class, $\sqrt{2}$ is irrational. A similar proof shows that $-\sqrt{2}$ is irrational (otherwise, if $-\sqrt{2} = \frac{a}{b}$, then $\sqrt{2} = \frac{-a}{b}$. However, $\sqrt{2} + (-\sqrt{2}) = 0$ which is rational.

9. Prove that there is an integer m such that $m^3 > 10^{10}$. Is your proof constructive or non-constructive?

Let $m = 2200$. Then $m^3 = (2200)^3 = 10,648,000,000 > 10^{10} = 10,000,000,000$.

This proof is constructive.

10. Prove that given any two rational numbers $p < q$, there is a rational number r with $p < r < q$.

Let p, q satisfy $p < q$. Let $r = \frac{p+q}{2}$.

Claim 1: $p < r$.

Since $p < q$, $q - p > 0$. Then $r = \frac{p+q}{2} = \frac{p+p+(q-p)}{2} = \frac{2p}{2} + \frac{q-p}{2} = p + \frac{q-p}{2} > p$.

Claim 2: $r < q$.

Since $q - p > 0$, $p - q < 0$. Then $r = \frac{p+q}{2} = \frac{q+q+(p-q)}{2} = \frac{2q}{2} + \frac{p-q}{2} = q + \frac{p-q}{2} < q$.

Claim 3: r is rational.

Since p and q are rational, $p = \frac{a}{b}$ and $q = \frac{c}{d}$ with a, b, c, d integers and $b, d \neq 0$.

Then $r = \frac{\frac{a}{b} + \frac{c}{d}}{2} = \frac{\frac{ad+bc}{bd}}{2} = \frac{ad+bc}{2bd}$. Since $ad + bc$ is an integer and $2bd$ is a non-zero integer, then r is rational.

Hence $r = \frac{p+q}{2}$ is a rational number satisfying $p < r < q$.

11. Prove that given a non-negative integer n , there is a unique non-negative integer m such that $m^2 \leq n < (m+1)^2$

Let m be $\lfloor \sqrt{n} \rfloor$. Then $m \leq \sqrt{n}$, so $m^2 \leq m \cdot \sqrt{n} \leq \sqrt{n} \cdot \sqrt{n} = n$.

Also, $m+1 > \sqrt{n}$, so $(m+1)^2 > (m+1) \cdot \sqrt{n} > \sqrt{n} \cdot \sqrt{n} = n$.

Hence $m^2 \leq n < (m+1)^2$.

To see uniqueness, suppose k is an integer so that $k^2 \leq n < (k+1)^2$. Then $k \leq \sqrt{n}$ (since otherwise, if $k > \sqrt{n}$, then $k^2 > k\sqrt{n} > \sqrt{n} \cdot \sqrt{n} = n^2$).

Similarly, $k+1 > \sqrt{n}$. So we have $k \leq \sqrt{n} < k+1$. Therefore, since both m and k are integers, we must have $m = k$.

12. Prove or disprove: Every non-negative integer can be written as the sum of at most 3 perfect squares.

This statement is false. Notice that 7 cannot be written as the sum of 3 squares:

$0 + 0 + 0 = 0$, $1 + 0 + 0 = 1$, $1 + 1 + 0 = 2$, $1 + 1 + 1 = 3$, $4 + 0 + 0 = 4$, $4 + 1 + 0 = 5$, $4 + 1 + 1 = 6$, $4 + 4 + 0 = 8$.

13. Which digits can occur as the final digit of the cube of an integer? Justify your answer.

Notice that $0^3 = 0$, $1^3 = 1$, $2^3 = 8$, $3^3 = 27$, $4^3 = 64$, $5^3 = 125$, $6^3 = 216$, $7^3 = 343$, $8^3 = 512$, and $9^3 = 729$.

From these examples, we see that any digit can occur as the final digit in the cube of an integer.

14. Use a paragraph to show that $(B - A) \cup (C - A) = (B \cup C) - A$.

Proof:

“ \subseteq ” Let $x \in (B - A) \cup (C - A)$. Then $x \in B - A$ or $x \in C - A$. We will consider these two cases separately.

Case 1: Suppose $x \in B - A$. Then, by definition of set difference, $x \in B$ and $x \notin A$. Since $x \in B$, then, by definition of set union, $x \in B \cup C$. Therefore, since $x \in B \cup C$ and $x \notin A$, we have $x \in (B \cup C) - A$.

Case 2: Suppose $x \in C - A$. As above, by definition of set difference, $x \in C$ and $x \notin A$. Since $x \in C$, then, by definition of set union, $x \in B \cup C$. Therefore, since $x \in B \cup C$ and $x \notin A$, we have $x \in (B \cup C) - A$.

Since these are the only possible cases, then $(B - A) \cup (C - A) \subseteq (B \cup C) - A$

“ \supseteq ” Let $x \in (B \cup C) - A$. Then, by definition of set difference, $x \in B \cup C$ and $x \notin A$. We will once again split into cases. Note that in both cases, we have $x \notin A$.

Case 1: Suppose $x \in B$. Then, since $x \notin A$, we have $x \in B - A$. Hence, by definition of set union, $x \in (B - A) \cup (C - A)$.

Case 2: Suppose $x \in C$. Then, since $x \notin A$, we have $x \in C - A$. Hence, by definition of set union, $x \in (B - A) \cup (C - A)$.

Since these are the only possible cases, $(B - A) \cup (C - A) \supseteq (B \cup C) - A$

Thus $(B - A) \cup (C - A) \subseteq (B \cup C) - A$. \square

15. Use a paragraph to show that $A \cup (A \cap B) = A$.

Proof:

“ \subseteq ” Let $x \in A \cup (A \cap B)$. Then, by definition of set union, $x \in A$ or $x \in A \cap B$. We will consider these two cases separately.

Case 1: Suppose $x \in A$. Then, the result is clear.

Case 2: Suppose $x \in A \cap B$. Then, by definition of set intersection, $x \in A$ and $x \in B$, so we again have the desired containment.

Hence $A \cup (A \cap B) \subseteq A$.

“ \supseteq ” Let $x \in A$. Then, by definition of set union, we must have $x \in A \cup (A \cap B)$. Hence $A \cup (A \cap B) \supseteq A$.

Thus $A \cup (A \cap B) = A$. \square

16. Use a paragraph proof to show that $A - B = A \cap \overline{B}$.

First, suppose $x \in A - B$. Then, by the definition of set subtraction, $x \in A$ and $x \notin B$. Thus $x \in \overline{B}$. Hence $x \in A \cap \overline{B}$. Therefore, $A - B \subseteq A \cap \overline{B}$.

Next, suppose $x \in A \cap \overline{B}$. Then $x \in A$ and $x \in \overline{B}$. Thus $x \notin B$. Hence $x \in A - B$. Therefore, $A \cap \overline{B} \subseteq A - B$.

Hence $A - B = A \cap \overline{B}$. \square